## A101 3次元可逆セル空間における自己増殖と形態形成 Self-reproduction and shape formation in three-dimensional reversible cellular space

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### Abstract

Reversible cellular automata (RCA) such as lattice gas models are used for modeling physical phenomena and they are also used for studying reversible computing processes. Due to power dissipation, it is said that nano-scaled computing devices should perform their computing processes in reversible manner. Although reversible computing processes can be computation-universal, it is quite difficult to place preferred initial configurations and advance computing. In this paper, we construct a three-dimensional self-inspective self-reproducing reversible cellular automaton. It can self-reproduce various shapes in three-dimensional reversible cellular space without dissipating any garbage and can be also used for placing preferred shapes in reversible environments.

Keywords: Cellular automata, Reversibility, Self-reproduction, Shape formation

## 1 Introduction

A reversible cellular automaton (RCA) is one of the reversible computing models. Its global function is injective and every configuration has at most one predecessor. Intuitively, it "remembers" the initial configuration and one can reconstruct its initial configuration from a configuration of any time. So reversible property is a strong constraint and one cannot generate nor extinct signals freely. Toffoli showed that there exist computationuniversal RCAs[11] and the BBM cellular automaton (BBMCA), was introduced by Margolus[6]. It is computation-universal and it has a direct relation with a physically reversible and conservative computing model (the Billiard Ball Model, BBM)[2]. The BBM has an important aspect that it is possible to compute any function without dissipation of balls as garbage, thus, it is possible to construct a computer that can computes with no energy dissipation in principle[1].

But the fact does not mean any computing process can be simulated effectively in reversible[8, 3]. It is very difficult to place a preferred initial configuration and start computing on reversible computing models. This problem is described as a restriction of generation and extinction of signals in RCAs, and synchronizing signals and distributing specific patterns on RCAs are also quite difficult. For example, von Neumann's (irreversible) self-reproducing cellular automaton[10] generates and erases many signals in its self-reproducing processes. This problem might have some relations with the difficulty of constructing three-dimensional functional solid nano-structures, because energy dissipation would be a big problem on them.

So we constructed a simple self-reproducing RCA based on a shape-encoding mechanism  $(SR_8)[9]$  to exhibit the possibility of such shape formation in a reversible space. In  $SR_8$ , self-reproduction can be performed without garbage in two-dimensional reversible cellular space.

In this paper, we extend  $SR_8$  into threedimensional reversible cellular space. Even if its cellular space is reversible, it can self-reproduce various three-dimensional patterns without garbage dissipation. In order to design an RCA we use a framework of partitioned cellular automaton (PCA). In the next section, first we define PCA.

## 2 Definitions

Partitioned cellular automaton (PCA)[7] is regarded as the subclass of standard cellular automaton. Each cell is partitioned into the equal number of parts to the neighborhood size and the information stored in each part is sent to only one of the

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neighboring cells (Fig. 1). In PCA, injectivity of global function is equivalent to injectivity of local function, thus a PCA is reversible if its local function is injective. Using PCA, we can construct a reversible CA with ease.

A deterministic two-dimensional partitioned cellular automaton (PCA) P is defined by

$$P = (\mathbb{Z}^2, (C, U, R, D, L), \varphi, (\#, \#, \#, \#, \#))$$

where  $\mathbb{Z}$  is the set of all integers, C, U, R, D, L are non-empty finite sets of center, up, right, down, left parts of each cell,  $\varphi : C \times D \times L \times U \times R \rightarrow$  $C \times U \times R \times D \times L$  is a local function (Fig.2), and  $(\#, \#, \#, \#, \#) \in C \times U \times R \times D \times L$  is a quiescent state which satisfies  $\varphi(\#, \#, \#, \#, \#) =$ (#, #, #, #, #).

A configuration over  $C \times U \times R \times D \times L$  is a mapping  $c: \mathbb{Z}^2 \to C \times U \times R \times D \times L$ . Let  $\operatorname{Conf}(C \times U \times R \times D \times L)$  denote the set of all configurations over  $C \times U \times R \times D \times L$ .

$$\begin{aligned} \operatorname{Conf}(C \times U \times R \times D \times L) = \\ \{c|c: \mathbb{Z}^2 \to C \times U \times R \times D \times L\} \end{aligned}$$

**Global** function

$$\Phi_A : \operatorname{Conf}(C \times U \times R \times D \times L) \rightarrow \operatorname{Conf}(C \times U \times R \times D \times L)$$

is defined by

$$egin{aligned} \Phi_A(c)(x) &= arphi( ext{CENTER}(c(x,y)), \ & ext{DOWN}(c(x,y+1)), \ & ext{LEFT}(c(x+1,y)), \ & ext{UP}(c(x,y-1)), \ & ext{RIGHT}(c(x-1,y))) \end{aligned}$$

where CENTER (UP, RIGHT, DOWN, LEFT, respectively) is the projection function which picks out the center (up, right, down, left) element of a quintuple in  $C \times U \times R \times D \times L$ . It has been proved that P is reversible iff  $\varphi$  is one-to-one[7].

P is called a rotation-symmetric (or isotropic) PCA iff (i) and (ii) hold.

(i) U = R = D = L.

- (ii)  $\forall (c, u, r, d, l), (c', u', r', d', l') \in C \times U^4$ :
- $egin{array}{lll} ext{if} & g(c,d,l,u,r) &= (c',u',r',d',l') & ext{then} \ g(c,r,d,l,u) &= (c',l',u',r',d'). \end{array}$

A deterministic three-dimensional partitioned cellular automaton (PCA)  $P_3$  is also defined by

$$P_{3} = (\mathbb{Z}^{3}, (C, U, R, D, L, F, B), \varphi_{3}, (\#, \#, \#, \#, \#, \#, \#, \#, \#))$$

Local function  $\varphi_3$  and global function are also defined in the same way as in the two-dimensional case.





Figure 2: Representations of a rule.



Figure 3: Domain and rage of local function in three-dimensional PCA.

## 3 Self-reproduction in a twodimensional RPCA

### 3.1 Definition of $SR_8$

In this section, we construct non-trivial selfreproducing structures can be constructible in a reversible cellular space. The idea of our model is based on Langton's sheathed Loop[5], and to achieve more flexibility we introduced a selfinspection method.

In the cellular space of  $SR_8[9]$ , encoding the shape of an object into a "gene" represented by a command sequence, copying the gene, and interpreting the gene to create an object, are all performed reversibly. By using these operations, various objects called Worms and Loops can reproduce themselves in a very simple manner.

The RPCA " $SR_8$ " is defined by

$$SR_8 = (\mathbb{Z}, (C, U, R, D, L), g, (\#, \#, \#, \#, \#)),$$
  

$$C = U = R = D = L = \{\#, *, +, -, A, B, C, D\}.$$

Hence, each of five parts of a cell has 8 states. The states A, B, C and D mainly act as signals that are used to compose "commands". The states \*, +, and - are used to control these signals. The local function g contains 765 rules. It is a one-to-one mapping and rotation-symmetric.

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wire  $(x_i, y_i \in \{A, B, C\})$ .

#### 3.2Signal transmission on a Wire

A wire is a configuration to transmit signals A, B, and C. Fig. 4 shows an example of a part of a simple (i.e., non-branching) wire.

A command is a signal sequence composed of two signals. There are six commands consisting of signals A, B and C as shown in Table 1. These commands are used for extending or branching a wire.

#### 3.3A Worm

A Worm is a simple wire with open ends that are called a head and a tail. It crawls in the reversible cellular space as shown in Fig. 5. Commands in Table 1 are decoded and executed at the head of a Worm. That is, the command AA extends the head straight, while the command AB (or AC, respectively) extends it leftward (rightward). On the other hand, at the tail cell, the shape of the Worm is "encoded" into an advance command. That is, if the tail of the Worm is straight (or left-turning, right-turning, respectively) in its form, the command AA (AB, AC) is generated. The tail then retracts by one cell.

#### $\mathbf{3.4}$ Self-reproduction of a Worm

By giving a branch command, any Worm can selfreproduce indefinitely provided that it neither cycles nor touches itself in the branching process.

#### 3.5Self-reproduction of a Loop

A Loop is a simple closed wire, thus has neither a head nor a tail as shown in Fig. 6.

If a Loop contains only advance or branch commands, they simply rotate in the Loop and selfreproduction does not occur. In order to make a Loop self-reproduce, commands in Table 2 are used.

By putting a command DB at an appropriate position, every Loop having only AA commands in all the other cells can self-reproduce in this way. When

Command		
First signal	Second signal	Operation
A	A	Advance the head forward
A	В	Advance the head leftward
A	С	Advance the head rightward
В	Α	Branch the wire in three ways
В	В	Branch the wire in two ways (making leftward branch)
В	С	Branch the wire in two ways (making rightward branch)

Figure 4: Signal transmission on a part of a simple Table 1: Six commands composed of A, B, and C.



Figure 5: Behavior of a Worm.

A+ A	A A +	
+ A A	A A+	A A +
Å A	+ A A	+ A + A

Figure 6: An example of a Loop.

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Com	mand	
First signal	Second signal	Operation
D	В	Create an arm
D	С	Encode the shape of a Loop

Table 2: Commands DB and DC.

DB reaches the bottom right corner, it starts making an "arm" and this corner become a transmitter of commands. First, all AA commands in the mother Loop are transmitted through the arm and generated DC commands encode whole shape of the mother Loop into command sequences simultaneously and these commands are transmitted after all static AA commands are transmitted. Finally DC commands reaches the bottom right corner and the arm is split from the mother Loop.

## 3.6 Controlling the position of daughter Loops in SR<sub>8</sub>

One of our main motivations is to place preferred initial patterns to a reversible cellular space. As mentioned above, a closed Loop has only AA commands. If AB or AC commands are placed in the Loop, generated position of the daughter Loop can be changed. But DB (create an arm command) advances the bottom side of a loop and the length of the Loop does not equal to the running length of the whole commands. Thus the embedded position of turning commands in the daughter Loop differ from the mother Loop. Although such a shifting of reading-frames of its command sequence is interesting phenomenon, it is difficult to control. So we modify  $SR_8$  for solving this timing problem.

Fig.7 is the process of modified version of  $SR_8$ . DB signal is not advance bottom side and the reproducing process starts from the bottom right corner as soon as the Worm reaches at this position. But created daughter Loop should be rotated in 90 degrees and Loops make collision after 4 generations because of this rotation. This collision can be avoidable by inserting direction commands into the mother Loop and this modification acts important roll in extending  $SR_8$  to three-dimensional one in the next section.

## 4 Self-reproduction in a threedimensional RPCA

## 4.1 Three-dimensional selfreproducing RPCA (SR<sub>9</sub>)

In this section, we extend  $SR_8$  into a threedimensional RPCA.

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Figure 7: Self-reproducing process of a Loop of modified  $SR_8$ .

A two-dimensional 5-neighbor PCA can be embedded directly into a three-dimensional 7-neighbor PCA. But due to the rotation-symmetric condition of  $SR_8$ , the Worm cannot know directions of right, left, up and down. In three-dimensional rotationsymmetric CA, up to 24 rotated rules are regarded as the same rule. So we introduce another glue state '=' for  $SR_8$  and combine two Worms whose command sequences are complementaly placed as presented in table 3. To construct 'true' threedimensional structures, a Worm in  $SR_9$  can twist its head in  $\pm 90$  degrees (Fig.8). This can be possible by employing ribbon of width 3 shaped Worms. As far as using  $SR_8$  command sequences in rotationsymmetric spaces, the length of the path should be kept equal and the width 2 ladder approach in the previous section is impossible. So we add a center wire and thus the three-dimensional selfreproducing RPCA " $SR_9$ " is defined by

$$SR_9 = (\mathbb{Z}^3, (C, U, R, D, L, F, B), g,$$



Figure 8: Four kind of turns in  $SR_9$ .

$$(\#, \#, \#, \#, \#, \#, \#, \#)),$$
  
 $C = U = R = D = L = F = B =$   
 $\{\#, *, +, -, =, A, B, C, D\}.$ 

Local rules are available via WWW. http://kelp.iec.hiroshima-u.ac.jp/ projects/rca/sr3d/

Although  $SR_9$  has 6886 rules, if rotated rules are regarded as equivalent, it become only 338 rules[4]. Fig.9 is a simple Worm in  $SR_9$ .



Figure 9: A simple Worm in  $SR_9$ .

Table 3: 'A'-Commands for width 3 shaped worm.

1	Com	nand		
wi	ne i	wig	23	
First Signal	Second Signal	Pirst Signal	Second Signal	Operations
Α	A	A	A	A dvance the head forward
A	B	A	С	Advance the bead left ward
A	С	A	B	Advance the head rightward
A	B	A	B	Start rotating (leftward)
A	C	A	C	Start rotating (rightward)

When both wires of a three-dimensional Worm have the same sequence 'AB AA AC' (or 'AC AA AB'), its head is twisted leftward (rightward). Using twisting commands, complex three-dimensional Worms and Loops such as Fig.10 are constructible. Although the existence of twisting commands in  $SR_9$ , its self-reproducing mechanism is completely the same as that of  $SR_8$ .

## 4.2 Controlling the position of daughter Loops in $SR_9$

When extending  $SR_8$  to  $SR_9$ , we use the modified version of  $SR_8$  discussed in section 3.6. So Loop positioning commands can also be inserted



Figure 10: Complex Worm and Loop in  $SR_9$ .

freely in  $SR_9$ . And this modification has an important meaning in the three-dimensional case because it makes possible to generate different topological shapes. Fig.11 is a chain formed from a single Loop. This shape-construction technique can be possible by the positioning a daughter Loop with a specific command sequences in the mother Loop.



Figure 11: A chain formed from a single Loop in  $SR_9$ .

# 5 Reflective behavior of loops in $SR_9$

In this section, we describe that  $SR_9$  has an ability to change shapes of daughter Loops even if its space is reversible.

Fig.12 shows an example of the generation of the different shape of daughter Loop. Although the mother Loop is square, its daughter Loop is rectangle. The length n of this square is 24 and it has the



Figure 12: Changing the shape of a daughter Loop in  $SR_9$ . (t=0,28,100,110)

of following positioning command sequence instead of  $AA^{24}$ .

AA<sup>5</sup>ABAA<sup>6</sup>AA<sup>3</sup>ABAAABAA<sup>6</sup> AA<sup>5</sup>ACAA<sup>6</sup>AA<sup>3</sup>ACAAACAA<sup>6</sup>

Once a duplication process starts, the shape is encoded as shape signals and they are transmitted along the arm after all positioning signal are transmitted. Thus the following command sequence (length 2n - 2 = 46) is generated when it is completely expanded into a Worm.

In normal self-reproducing process of  $SR_9$ , daughter Loops are constructed after all positioning commands are used for advancing the constructing arm. But in this case, the daughter Loop is closed soon after its arm are split from the mother Loop. Thus the positioning signals and the shape signals are exchanged.

The next example is more complex one. The diagram of its shape is depicted in Fig.13(a) and it has the following command sequence.

(AA AB AA AC AA AC AA AB AA $^4$ AA $^6$ AB AA AC AB AA AB AA AC AA AC AA AC AA AB AA AB AA $^7$ AB AA) (AA AB AA AC AA AC AA AB AA $^4$ AA $^6$ AC AA AB AC AA AC AA AB AA AC AA AC AA $^7$ AB AA)
·
positioning signals
AA ACAA ABAA ABAA ACAA <sup>2</sup> ABAA <sup>7</sup> ABAA <sup>2</sup> ABAA ACAA ACAA ABAA <sup>2</sup> ABAA
AA AB AA AC AA AC AA AB AA <sup>2</sup> AC AA <sup>7</sup> AC AA <sup>2</sup> AC AA AB AA AB AA AC AA <sup>2</sup> AC AA <sup>7</sup>



Each side of this square has specific patterns and they are changed after self-reproduction. Fig.13(b) is generated from Fig.13(a) when start symbol D is provided at the timing of the command sequence shown above. Underlined part is used to form its shape. This shape/positioning command sequence can generate four different Loops in turn.



Figure 13: The diagram of changing the shape of a daughter Loop (Period 4) in  $SR_9$ .

## 6 Conclusion

In this paper, we extend our two-dimensional self-reproducing reversible PCA  $SR_8$  into a threedimensional reversible PCA and show its various behaviors. The features of this three-dimensional reversible "turtle graphics" are derived from its selfinspective mechanism. Data as shapes and programs as command sequences are represented in the same manner.

The self-reproducing processes are hard to describe on a paper. They can be seen as QuickTime Movies at the following addresses via WWW.  $SR_9$ : http://kelp.iec.hiroshima-u.ac.jp/projects/rca/sr3d/

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