

2-A-6

AN OPTIMIZATION-ORIENTED SIMULATION-BASED JOB SHOP SCHEDULING
METHOD WITH FOUR PARAMETERS USING PATTERN SEARCH

Masahiro Arakawa
Department of Systems
Management Engineering
Kansai University
Suita, Osaka 564-8680
Japan
arakawa@iecs.kansai-u.ac.jp

Masahiko Fuyuki
Department of Systems
Management Engineering
Kansai University
Suita, Osaka 564-8680
Japan
fuyuki@iecs.kansai-u.ac.jp

Ichiro Inoue
Faculty of Management
Kyoto Sangyo University
Kita-ku, Kyoto 603-8047
Japan
inoue@cc.kyoto-su.ac.jp

ABSTRACT

Aiming at the elimination of tardy jobs in a job shop production schedule, an optimization-oriented simulation-based scheduling (OSBS) method incorporating capacity adjustment function is proposed. In order to determine the pertinent additional capacities and to control job allocations simultaneously the proposed method incorporates the parameter-space search improvement (PSSI) method into the scheduling procedure. In previous papers, we have introduced four parameters; two of them are used to control the upper limit to the additional capacity and the balance of the capacity distribution among machines, while the others are used to control the job allocation procedure. We found that a 'direct' optimization procedure which uses the enumeration method produces an optimal solution with practical significance, but it takes too much computation time for a practical use. In this paper, we propose a new method which adopts a pattern search method in the schedule generation procedure to obtain an approximate optimal solution. It is found that the computation time becomes short enough for a practical use. Moreover, the extension of the parameter domain yields an approximate optimal solution which is better than the optimal solution obtained by the 'direct' optimization.

INTRODUCTION

In a make-to-order production, the most important requirement for production scheduling is to fulfill the customer due-date for each job. When tardy jobs are involved in a production schedule, an improved schedule is sought for to eliminate tardy jobs under the given available capacities. When the occurrence of the tardy job is judged to be inevitable, the reduction of the number of tardy jobs is attempted by increasing resource capacities within the allowances granted by the upper-level planning department. However, in job shop production scheduling, it is generally difficult to decide the capacities when, to which machine and what amount, to add so as to minimize of the number of tardy jobs after suppressing the added amount of the capacity.

In order to determine the pertinent additional capacities and to control job allocations simultaneously we incorporate

the parameter-space-search-improvement method (named the PSSI method) (Fuyuki *et al.* 1998) into the scheduling procedure. The PSSI method is a framework for finding the best solution in a simulation method. It introduces a very few number of parameters which can systematically manipulate the relevant variables and the best solution is sought for on the parameter space spanned by the introduced parameters. In previous papers, we have introduced four parameters; two of them a and b are used to control the upper limit to the additional capacity and the balance of the capacity distribution among machines (Arakawa *et al.* 2001a). The best capacity addition plan is sought for on the parameter space $a \times b$, while the others c_1 and c_2 are used to control the job allocation procedure (Fuyuki *et al.* 1999) and the best schedule with respect to due-date related criterion is sought for on the parameter space $c_1 \times c_2$. A 'direct' optimization procedures on the four-dimensional direct-product space which we named the OSBS/4 method is developed by merging the procedures separately treated on the two dimensional spaces (Arakawa *et al.* 2001b). By using scheduling data obtained from a practical large-scale system, the performance of the procedure was investigated. It is found that the number of tardy jobs and the added amount of the capacity are significantly reduced, but the computation time to obtain the optimal solution by enumeration is too long for practical application.

In order to shorten the computation time we propose in this paper a new method which adopts a local search method in the schedule generation procedure, and examine the performance of the propose method.

AN OPTIMIZATION-ORIENTED SIMULATION-BASED SCHEDULING METHOD

As for scheduling aiming at the reduction of job tardiness, a certain type of the Backward/Forward Hybrid Simulation method (named BFHS/type C) is known to show excellent performance over the conventional forward simulation method with sophisticated dispatching rules (Fuyuki and Inoue 1995). In the BFHS/type C method, the due-date related information obtained from the first step backward simulation is utilized in the second step forward simulation, i. e., the order of the

starting time for each job at each work center estimated in the backward simulation is used in the forward simulation to create a job priority order for the concerned jobs at each work center. We adopt this BFHS/type C method as a schedule generation method in the scheduling procedures in this section.

Parameters to control job allocations

Since the set of the parameters (c_1, c_2) shows the best performance with respect to tardiness related criteria such as the average tardiness, the maximum tardiness and the number of tardy jobs (Fuyuki *et al.* 1999), we adopt this parameter set to control the job allocation process in an optimization procedure.

This parameter set was introduced with the intention to control the job allocation process by utilizing the characteristics of the initially generated scheduling result. Focusing on the due-date lateness (= the job completion time – the due-date) and the accumulated waiting time of each job in the initial schedule, we intended to adjust the slack time for each job at every occasion when it is estimated for calculating the job priorities. In order to reduce the number of the degrees of freedom, we introduce the two parameters and decide the adjustment quantity by the following equation

$$\Delta_i = c_1 L_i + c_2 W_i, \quad (1)$$

where Δ_i is the adjustment quantity for the slack time for a job i , and the coefficients c_1 and c_2 are the parameters. By systematically changing the parameter values, we can control the job allocation process.

Parameters to control capacity adjustments

When we limit the capacity adjustment to the supplement of capacities, the items to be determined in the capacity adjustment control procedure are the capacities when, to which machine and what amount, to add. In order to make a pertinent decision on the capacity supplements, we introduced two parameters a and b which can systematically manipulate the relevant decision variables (Arakawa *et al.* 2000). The parameter a which takes the range $0 \leq a \leq 1$ controls the upper limit to the additional capacity and the parameter b which takes the range $-1 \leq b \leq 1$ the balance of the capacity distribution among machines. Detail account of the determination of the “initially added capacity” for a machine per day is given in Arakawa *et al.* (2000).

Schedule generation procedure

A ‘direct’ optimization procedure to seek for an optimal solution on the direct product space of the four parameters a, b, c_1 and c_2 was developed by Arakawa *et al.* (2001b), where the enumeration method is adopted to select an optimal solution. In order to reduce the computation time, we adopt one of a direct search method termed the pattern search (Hooke and Jeeves, 1961). Since the search method is adopted, the solution to be found can be an approximate optimal solution. However, the extension of the original parameter domain by the adoption

of finer step sizes (grid spacing) and by the removal of the parameter boundaries yields an approximate optimal solution which can be better than the optimal solution defined on the original parameter domain.

The proposed scheduling procedure which we shall name the optimization-oriented simulation-based scheduling method with four parameters using a pattern search method (*the OSBS/4/P method* in abbreviation) starts with the initial schedule generation and calculation of necessary quantities followed by a pattern search procedure with modifications.

The pattern search procedure is described in Appendix B by Hooke and Jeeves (1961) in a general form. The procedure is applied to the current four-dimensional space problem: a space point is specified by the coordinates a, b, c_1, c_2 ; the ‘function’ is the procedure at a given space point to calculate the number of tardy jobs in a schedule and to decide the total additional capacity; and the ‘functional value’ corresponds to the number of tardy jobs as well as the total additional capacity.

We introduce the following modifications into the original procedure:

- (1) In the original procedure, not all of the possible exploratory points are visited, and the new coordinate after a success move is retained and used at later steps. We modify this procedure to check all possible values obtained in exploratory moves, since ‘functional value’ is found to change discontinuously on the parameter space.
- (2) When the reduced step size becomes smaller than the “minimum” step size δ , reset the base point by a move to a randomly chosen point and set the step size the initial value, and restart the pattern search as long as the accumulated number of the base point random jump is smaller than the preset limit. This modification is expected to allow a search not to be trapped in a local minimum on the parameter space.

During the search procedure we refer to the solution and the ‘functional value’ so far attained, which we shall call the reference solution and the reference value, respectively.

The OSBS/4/P method is described as follows:

Main procedure to search for an approximate optimal solution

Step 1. Search condition setup

Set the following values related to the search condition: the initial step size Δ_0 , the ‘minimum’ step size δ , the initial position p_0 of the base point, the maximum number M of the base point random jump along with the stochastic function which gives a size of the random jump.

Step 2. Initial schedule generation and calculation of necessary quantities

Generate an ‘initial schedule’ subject to the initial planning condition by the BFHS/type C method, and calculate the quantities which will become necessary in the subprocedure: the machine utilization, the due-date lateness L_i and the accumulated waiting time W_i for every job i .

Step 3. Initial value assignments

Assign the initial value p_0 to the base point p , the initial

value Δ_0 to the step size Δ , the initial schedule and the initial planning condition to the reference solution, and the number of tardy jobs in the initial schedule and the value zero of the total additional capacity to the reference value.

Step 4. Exploratory points arrangement and functional values estimation

Step 4a. Exploratory points arrangement

Arrange two exploratory points for each coordinate which locate both side of the base point p with the distance Δ . Since we have four coordinates, eight exploratory points are configured around the base point. When the coordinate values for a or b exceed their own boundary, put the exploratory point at each boundary point.

Step 4b. Solution generation and reference solution update

Repeat the following (1) and (2) for all eight exploratory points.

- (1) At each exploratory point x , execute the subprocedure $S(x)$ to obtain an exploratory solution and the functional value.
- (2) Compare the functional value with the reference value. If the functional value is superior to the reference value, update the reference solution and the reference value by the exploratory solution and the functional value, respectively.

Step 5. Base point and step size update

According to whether or not the reference solution is updated at Step 4b, take the following actions:

- (1) If the reference solution was updated, by using the exploratory point r where this reference solution is obtained, move the current base point p to the new base point $2r-p$, reset the step size Δ to the initial value Δ_0 , and go to Step 4.
- (2) Otherwise, move the current base point p to the point x where the reference solution is obtained and replace the value of the step size Δ by one-half of it. If new step size is less than the minimum step size δ , go to Step 6, or else go to Step 3.

Step 6. Search termination or a base point jump

If the accumulated number of the base point jump is equal to the maximum number M , set the current reference solution to the approximate optimal solution, and terminate the search procedure.

Otherwise, reset the step size Δ to the initial value Δ_0 , move the current base point p to the point $p+rand$ where $rand$ means the value given by the random number with the normal distribution, and go to Step 4.

Subprocedure $S(x)$

Step s1. Capacity allocation

Determine the amount of additional capacities and the dates to add for each machine in terms of the coordinate values on the exploratory point x and distribute them to each machine. We shall call these capacity distributed at this step "initially added capacity."

Step s2. Schedule generation by the BFHS/type C method with slack time modifications

Allocate the jobs to the machines by the BFHS/type C method subject to the planning condition that extends the initial planning condition to include the initially added capacity. In the backward simulation process of the BFHS/type C method, the slack time modifications specified by the parameters c_1 and c_2 are taken into account.

Step s3. Elimination of redundant capacity

On the basis of the schedule generated at Step s2, eliminate the "unused time period" from the initially set over time period (i. e., the added capacity on that day). Here we call the period that no job is allocated until the end of that day the "unused time period". Moreover, if it is found that the added capacity on each machine on a particular day is not efficiently used, cancel the capacity addition on that machine on that day. We regard in this paper the capacity addition inefficient if the time interval between the end of the normal available time period and the onset time of a job allocated in the added capacity time period is larger than 30 minutes.

Step s4. Schedule regeneration by the BFHS/type C method with slack time modifications

Reallocate the jobs to the machines by BFHS/type C method taking account of the capacity alteration at Step s3 and of the slack time modification.

Step s5. Elimination of "unused" capacity

On the basis of the schedule generated at Step s4, if an "unused time period" is again found, eliminate it from the modified over time period. At this step, the total additional capacity to all machines is fixed for a particular point on the parameter space.

Step s6. End of the subprocedure

Return the schedule generated at Step s4 and the amount of the total additional capacity calculated at Step s5.

When the estimated functional value gives the same number of tardy jobs so far attained at Step 4b(2), we compare the total additional capacities and if the present value is below that at another one with the same number of tardy jobs, we regard the functional value 'superior'.

Because of the modification (1), there is possibility to select the space point once visited, we store on the computer memory every calculated result as the attribute of a space point, and reuse it to speed up the computation.

EVALUATION OF THE PROPOSED METHOD

We examine the performance of the proposed method focusing on the time required to obtain the approximate optimal solution which has significance in a practical sense.

Since the OSBS/4 method which adopts the enumeration method shows excellent results on the scheduling data sets obtained from a large-scale system, we use the same data sets and compare the performance of the proposed method with that of the OSBS/4 method. The 48 data sets are extracted from the actual data in the currently working scheduling system obtained

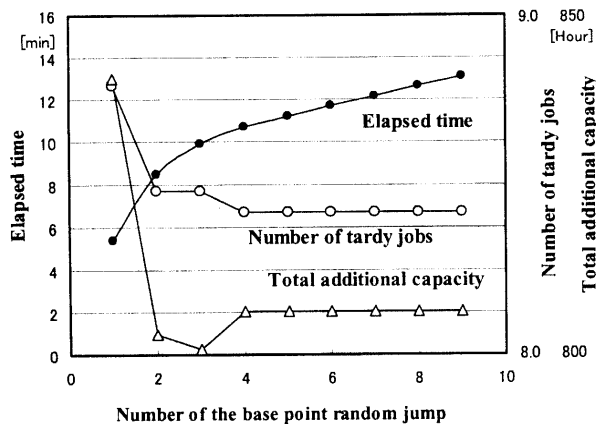


Figure 1. DEPENDENCE ON THE NUMBER OF THE BASE POINT RANDOM JUMP.

in the past 48 months. The characteristics of the actual data is given in Arakawa *et al.* (2001b).

The elapsed time required to obtain the approximate optimal solution by the proposed method is determined by the termination condition at *Step 6*, i.e., by the maximum number M of the base point random jump. The elapsed time T is directly governed by M , but also indirectly by the 'minimum' step size δ , because the number of the exploratory moves depend on δ at *Step 5*.

The values for M and δ should be chosen so that the practically significant solution can be derived by the proposed procedure. As an example for the practically significant solution, we cite the optimal solutions obtained by the OSBS/4 method which give the average values for the number of tardy jobs 8.7, and that for the total added capacity 822.0 (hour). The value of δ is to be smaller than the values 0.1 and 0.2 adopted in the OSBS/4 method. We shall choose the value of $\delta = 0.0125$.

The initial values for the initial step size Δ_0 and the initial position p_0 of the base point as well as the distribution of the stochastic function which gives a size of the random jump also affect the quality of the approximate optimal solution. The initial base point position is considered to be less important, because its effect may be compensated by the increase of M . On the other hand, the effects caused by the change of the initial step size Δ_0 and the distribution of the stochastic function are considered to be mutually related by themselves. We can investigate, for example, the dependence of the solution quality on the size of Δ_0 keeping the distribution constant.

First, let us investigate how the elapsed time T depends on M . Figure 1 shows the elapsed time measured on the PC (CPU: Pentium III, 1GHz) when the M -th base point random jump occurred. The increments of the elapsed time decrease as M increases. This is due to the device to save the computation time stated at the last paragraph of the previous section. The

number of tardy jobs and the total additional capacity shown on the same figure indicate that after the first base point jump the significant solution can be almost attained and a choice $M=3$ can be sufficient in practice. The elapsed time at $M=3$ is about 10 minutes which is by far smaller than 13 hours of the time required to obtain the optimal solution in the OSBS/4 method, and the proposed method can be used for an actual problem.

The results of Fig. 1 are obtained by setting the initial step size Δ_0 to 0.8, the value of which was found to be the best value. We have investigated the dependence of the quality of the solution on the size of Δ_0 and it was found that the number of tardy jobs decreases by about 2 jobs as Δ_0 increases from 0.2 to 0.8, and beyond $\Delta_0 = 0.8$ it slightly increases.

In conclusion, the computation time required in the propose method dramatically reduced to the level applicable to the practical use while keeping the number of tardy jobs as well as the total added capacity as the same or even slightly better than those of the OSBS/4 method.

A comparison of performances of the proposed search method with other search methods such as the original pattern search method and simulated annealing method is not presented in this paper, and it will be reported elsewhere.

REFERENCES

- ARAKAWA, M., FUYUKI, M., NAKANISHI, H., and INOUE, I. 2001a, A simulation-based method for capacity adjustments in job shop production scheduling. *Journal of Japan Industrial Management Association*, **51**, 603-612.
- ARAKAWA, M., FUYUKI, and INOUE, I. 2001b, An optimization method for simulation-based job shop scheduling incorporating capacity adjustment function", *16th International Conference on Production Research*, Prague, Czech, July 29 – August 3, 2001, CD-ROM, ISBN 80-02-01438-3, A7.4.
- FUYUKI, M., and INOUE, I. 1995, Due-date-conformance oriented production scheduling in a make-to-order production on the basis of backward/forward hybrid simulation. *Journal of Japan Industrial Management Association*, **46**, 144-151.
- FUYUKI, M., ARAKAWA, M., FURUICHI, Y., and INOUE, I., 1998, An improvement method for due-date-conformance oriented production schedule generated by simulation-based scheduling system. *Journal of Japan Industrial Management Association*, **48**, 370-377.
- FUYUKI, M., ARAKAWA, M., SUGIHARA, T., MATANO, K., and INOUE, I., 1999, A parameter space setting method in the improvement phase of simulation-based production scheduling. *Transactions of the Institute of Systems, Control and Information Engineers*, **12**, 63-65.
- HOOKE, R., and JEEVES T.A., 1961, Direct search solution of numerical and statistical problems. *Journal of Association of Computing Machinery*, **8**, 212-229.